

# Formal Analysis: Lecture 4

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# Extensive Games vs. Strategic Games

## **Strategic Games:**

We assume that players make decisions simultaneously (or without observing each other's decisions).

We assume that players can commit to a binding plan of action at the beginning of the game.

Appropriate for static situations.

## **Extensive Games:**

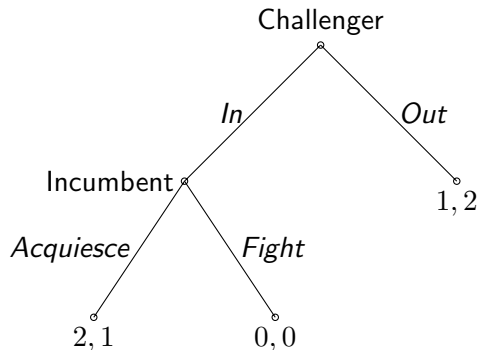
We assume that players make decisions sequentially.

We assume that some players can make up their mind after observing the previous actions of other players.

Appropriate for dynamic situations.

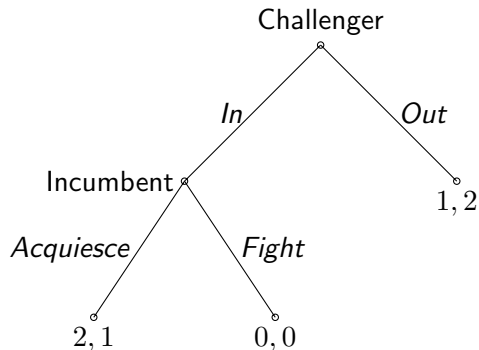
# Extensive Form

- The graphical representation is called a game tree



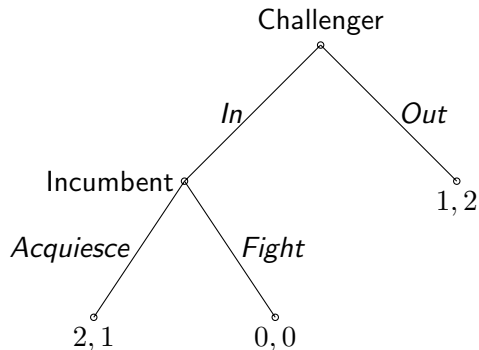
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- Extensive form games have four components:



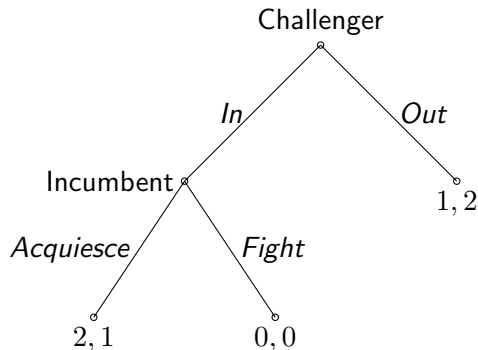
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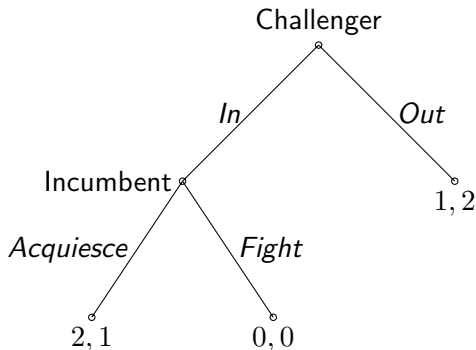
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  - ▶ a set of players
  - ▶ a set of sequences (terminal histories)



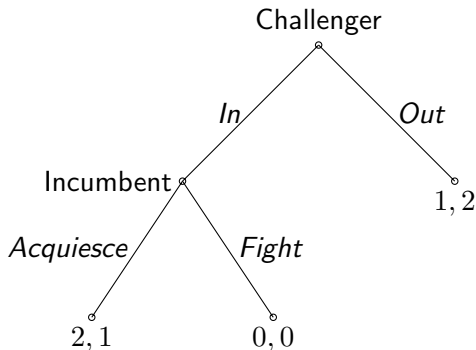
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  - ▶ a set of players
  - ▶ a set of sequences (terminal histories)
  - ▶ a player function that assigns a player to every proper subsequence of some terminal history
  - ▶ preferences over the set of terminal histories



# Extensive Form

- In this class we will study games with a finite horizon, a finite number of actions and perfect information.

**Finite horizon:** Terminal histories are not infinitely long. The game has an end.

**Finite number of actions:** Each player has a limited number of actions to choose from. There is a finite number of terminal histories.

A game with a finite horizon and a finite number of actions is called a finite game.

**Perfect information:** Each player knows all the actions taken previously (or at least knows what the exact situation is at each point in time). Each player always moves alone at each stage.

- In future classes we will study games with infinite horizons, infinite number of actions and/or imperfect information

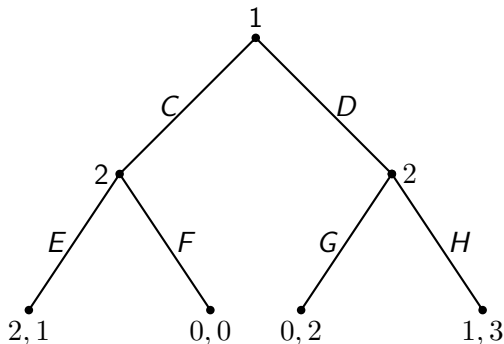
# Extensive Form

## Solving extensive form games

**Strategy:** A function that assigns an action to each *history* in which it is the player's turn

**Important:** A strategy must specify what a player will do in every contingency!

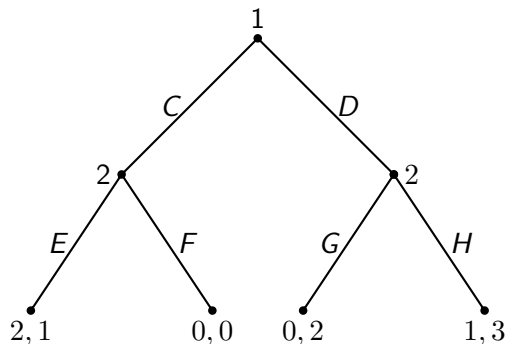
**In other words:** A strategy is a complete plan of action for every situation in which the player might be called upon to act.



# Extensive Form

What are all the possible strategies of player 2?

	Action assigned to history C	Action assigned to history D
Strategy 1	E	G
Strategy 2	E	H
Strategy 3	F	G
Strategy 4	F	H



# Nash Equilibrium

- We refer to a profile of strategies as  $s$ .
- The outcome of game is determined by the actions taken along the path of play specified by the strategies in  $s$ .
- We refer to outcomes as  $O(s)$

## Definition (Nash Equilibrium)

The strategy profile  $s^*$  is a NE if for all players and every strategy  $r_i$  the terminal history  $O(s^*)$  is at least as good as those generated by any deviation.

$$u_i(O(s^*)) \geq u_i(O(r_i, s_{-i}^*)), \forall r_i \text{ of player } i$$

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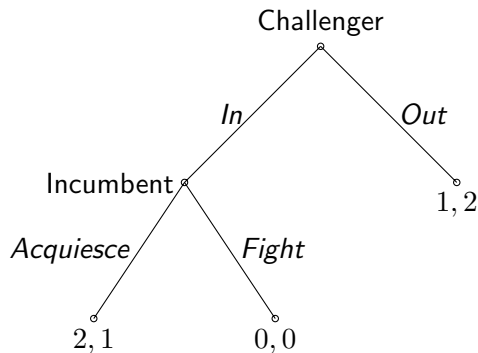
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## Definition

$$u_i(\alpha^*) \geq u_i(\alpha_i, \alpha_{-i}^*), \forall \alpha_i \in A_i$$

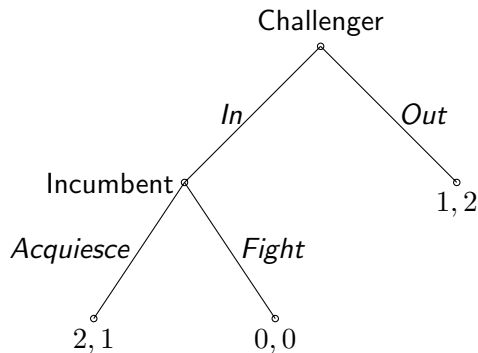
# Nash Equilibrium

- Consider the Incumbent vs. Challenger game again



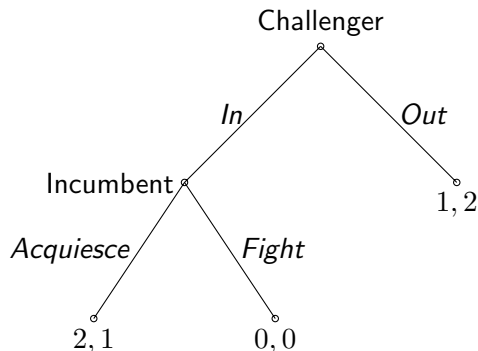
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- To find the the NE we need three components:



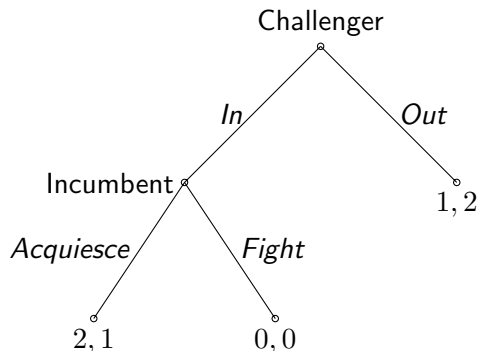
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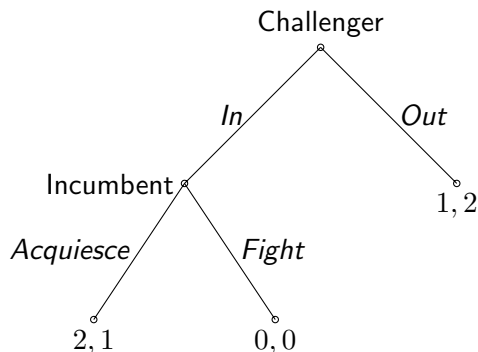
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- Consider the Incumbent vs. Challenger game again
- To find the the NE we need three components:
  - ▶ the set of players
  - ▶ the set of strategies (the actions chosen in each node)
  - ▶ preferences over the set of possible outcomes



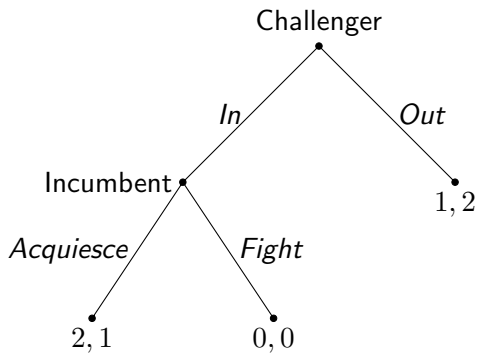
# Nash Equilibrium

- Thus, one way to find the NE is to construct a corresponding strategic form game
- Challenger vs. Incumbent game in strategic form:

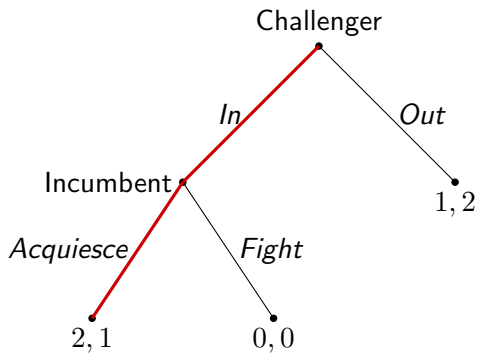
		Incumbent	
		<i>Acquiesce</i>	<i>Fight</i>
Challenger	<i>In</i>	2, 1	0, 0
	<i>Out</i>	1, 2	1, 2

- This game has two Nash equilibria: (*In*, *Acquiesce*) and (*Out*, *Fight*)

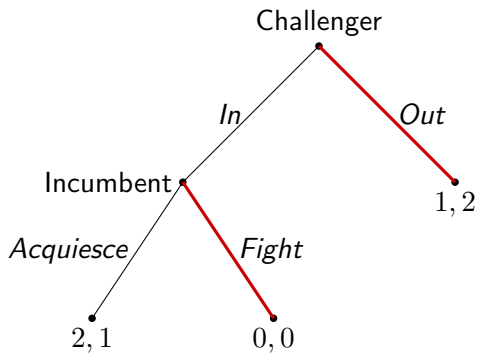
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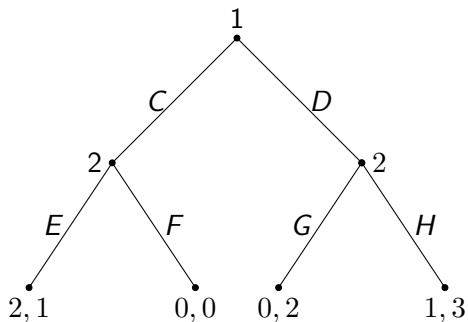
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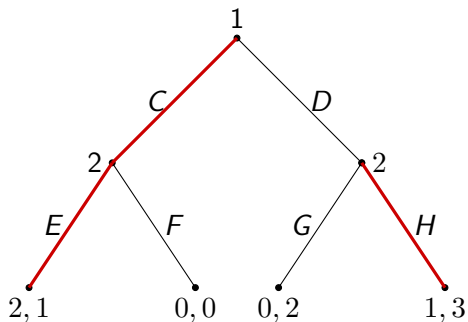
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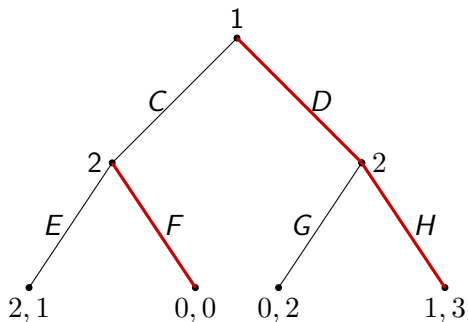
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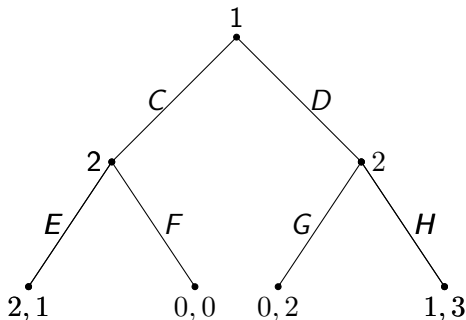


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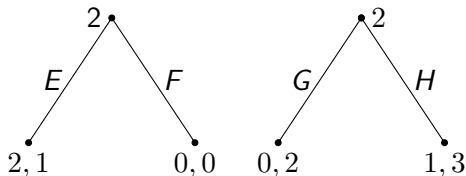
# Subgame Perfect Equilibrium

- So the concept of Nash equilibrium makes "too many" predictions. Can we find a different concept that does not predict non-credible threats?
- We start by defining the concept of a **subgame**. A subgame  $\Gamma(h)$  is the game following history  $h$ . The whole game is a subgame of itself. The other subgames are called **proper subgames**.



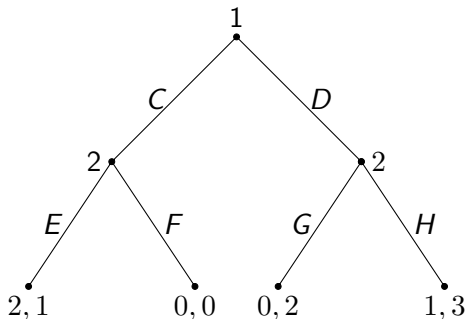
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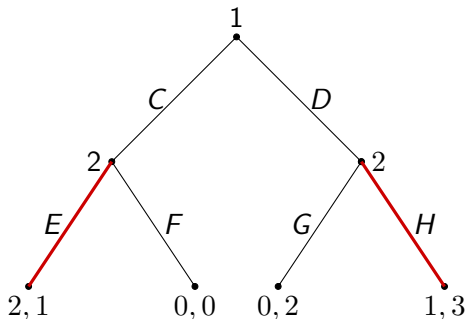
# Subgame Perfect Equilibrium

- A more refined equilibrium concept is that of **Subgame Perfect Equilibrium (SPE)**
- A player's actions must be optimal in every subgame of the game.
- Every SPE is a NE but not vice versa.
- A SPE is a strategy profile that induces a NE in every subgame.



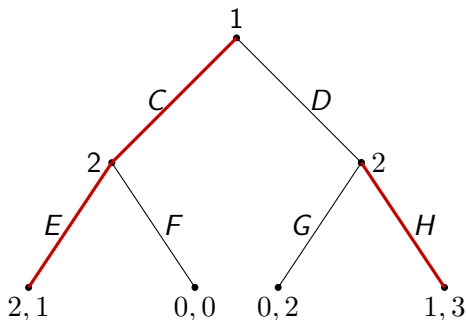
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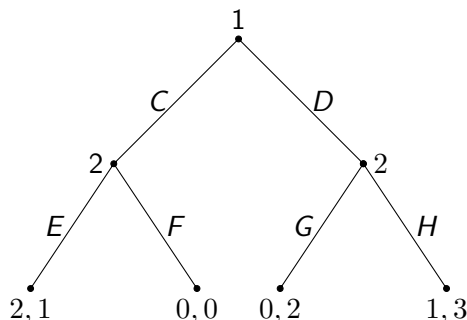
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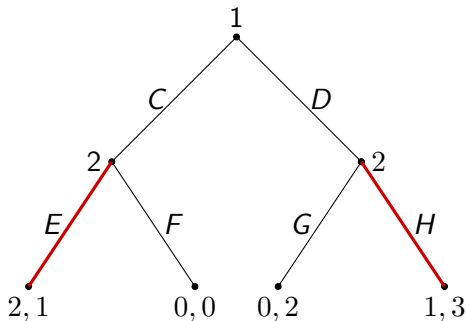
# Backwards induction in finite horizon games

- How do we find the SPE?
- Start at the bottom of the tree (the subgames with least length).
- Find the optimal action for each player at that stage of the game
- Construct the **reduced game** that takes the subsequent optimal actions as given. Apply the same procedure to the bottom of that new tree.
- This process is called **backwards induction**.



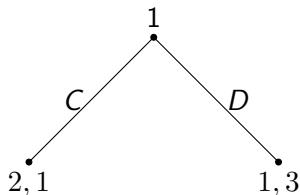
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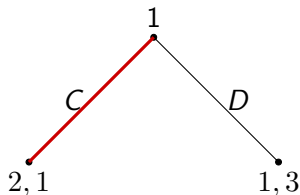
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# Conclusions

- The concept of SPE allowed us to eliminate some undesirable predictions. It is a refinement of NE.
- But do SPE always exist? Are they difficult to find?

## Theorem (Existence of subgame perfect equilibria)

*In every extensive game with a finite horizon, a finite number of actions and perfect information there exists at least one SPE.*

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## Theorem (Existence of subgame perfect equilibria)

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## Theorem

*If an extensive game has a finite horizon and perfect information then its subgame perfect equilibria can be found by backwards induction.*